

Generalization Correctness

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Problem Statement

- Given
 - System Model
 - Constraint
 - Solution provided by Constraint Solver
- Generate a Generalization
 - Convert a single solution into a set of solutions
 - Express Result Concisely
 - Usually Generalization != Constraint
 - Result is Inexact



Generalization



Generalization Illustration

- Computed via Symbolic Simulation
 - System Model + Constraint
 - Original Solution

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- Simulation is Approximate (Lossy)
 - Representational constraints

X = T Y = F Z = T Constraint = T

Generalization

Model

- Is the Generalization Correct?
 - Formalize Correctness
 - Articulate Generalization Rules
 - Prove Rules Satisfy Correctness



Generalization Correctness Statements

- Top Level Correctness Statement
 - Generalization Contains Original Solution
 - Generalization is a Subset of Original Constraint
- Invariants

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- Can be enforced incrementally
 - During Symbolic Simulation
- Reduce to Correctness when applied to top level constraint



- Correctness Invariants
 - 1. Evaluating Solution on Generalization must be the same as Evaluating Solution on original expression
 - 2. An input whose evaluation differs from that of the solution on the original expression must also differ on the Generalization





Generalization Rules

- Generalizing Boolean Expressions
 - AND, OR, NOT, ID
- One Choice:
 - Drop Terms or Not?
- Visualization
 - State Space
 - Original Solution is one Point
 - Organized as Truth Table w/to A,B
- Consider rules for Generalizing AND
 - OR follows from De Morgan's





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Rule #1: (AND F F)



- Correctness Invariants
 - 1. Evaluating Solution on Generalization must be the same as Evaluating Solution on original expression
 - An input whose evaluation differs from that of the solution on the original expression must also differ on the Generalization
- Generalization Rule #1
 - If both expressions evaluate to False, we can either keep both or keep just one



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Rule #2: (AND T T)



- Correctness Invariants
 - 1. Evaluating Solution on Generalization must be the same as Evaluating Solution on original expression
 - 2. An input whose evaluation differs from that of the solution on the original expression must also differ on the Generalization
- Generalization Rule #2
 - If both expressions evaluate to True, then we must keep both



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Rule #3: (AND T F)



- Correctness Invariants
 - 1. Evaluating Solution on Generalization must be the same as Evaluating Solution on original expression
 - An input whose evaluation differs from that of the solution on the original expression must also differ on the Generalization
- Generalization Rule #3
 - If the expressions evaluate to different values, we can either keep both or keep just the False expression





ACL2 Model

- Defined an expression evaluator
 - Expression and variable binding
 - AND, OR, NOT, IDs
- Used encapsulation to characterize 3 Generalization rules for AND
 - Choice is .. pragmatic
- Defined a depth-first generalizer
 - Returns a "generalized" expression
 - NOT, ID performs no simplification
 - Encapsulated function generalizes AND expressions
 - De Morgan's rule to simplify OR
- Formalized Correctness Invariants
- Proved that generalizer satisfied invariants





Expression Evaluator

```
(defun eval-expr (expr env)
(case-match expr
  (('and x y))
   (let ((x (eval-expr x env))
         (y (eval-expr y env)))
     (and x y)))
  (('or x y)
   (let ((x (eval-expr x env))
         (y (eval-expr y env)))
     (or x y)))
  (('not x)
   (let ((x (eval-expr x env)))
     (not x)))
  (('id n)
   (nth n env))
  (& expr)))
```





Generalizer Formalization

```
(defun gen-expr (expr sln)
(case-match expr
  (('and x y))
   (let ((genx (gen-expr x sln)))
         (geny (gen-expr y sln)))
                                         Applies 'and'
     (gen-and genx geny sln))) ←
                                             Rules
  (('or x y)
   (let ((genx (gen-expr x sln)))
         (geny (gen-expr y sln)))
     (gen-or genx geny sln)))
  (('not x)
   (let ((genx (gen-expr x sln)))
     (not-expr genx)))
  (& expr)))
```





Invariant Proofs



- Correctness Invariants
 - 1. Evaluating Solution on Generalization must be the same as Evaluating Solution on original expression
 - 2. An input whose evaluation differs from that of the solution on the original expression must also differ on the Generalization





Invariant Proofs



- Correctness Invariants
 - 1. Evaluating Solution on Generalization must be the same as Evaluating Solution on original expression
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Rule #3: (AND T F)



- Generalization Performed Depth-First
 - Solution space may get smaller (per correctness statement)
 - Predicate boundaries move closer to original solution
- Generalization Rule #3
 - If the expressions evaluate to different values, we may keep only the False expression





Conclusion

- We assumed that "Doing Nothing" was conservative
 - If you never change the expression, it trivially satisfies correctness
- We were wrong !
- It is easy to make these kinds of mistakes
 - ACL2 can help during algorithmic development
- Accomplishments
 - Formalized a notion of correctness for Generalization
 - Formalized rules for Generalization
 - Proved Generalization procedure
 - Corrected an error in our original Generalization rules